Introduction to

Algorithm Design and Analysis

[01] Model of Computation

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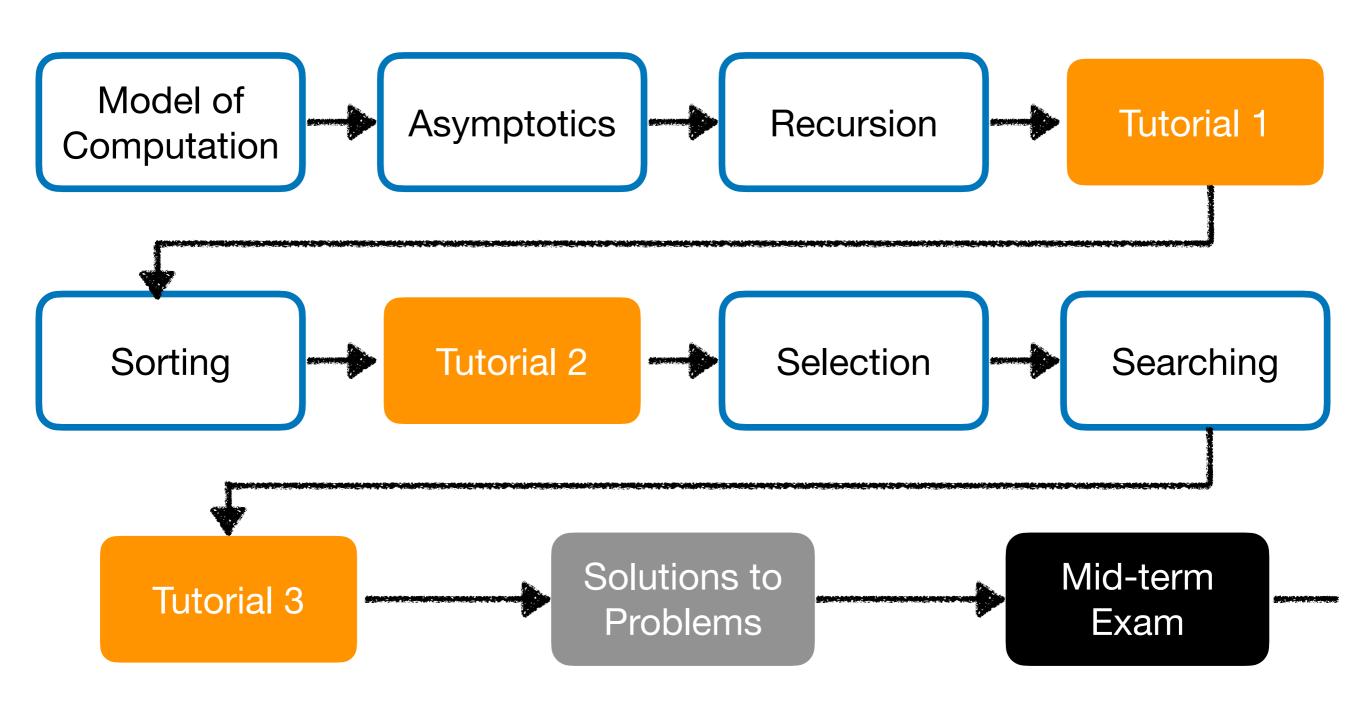
Course Information

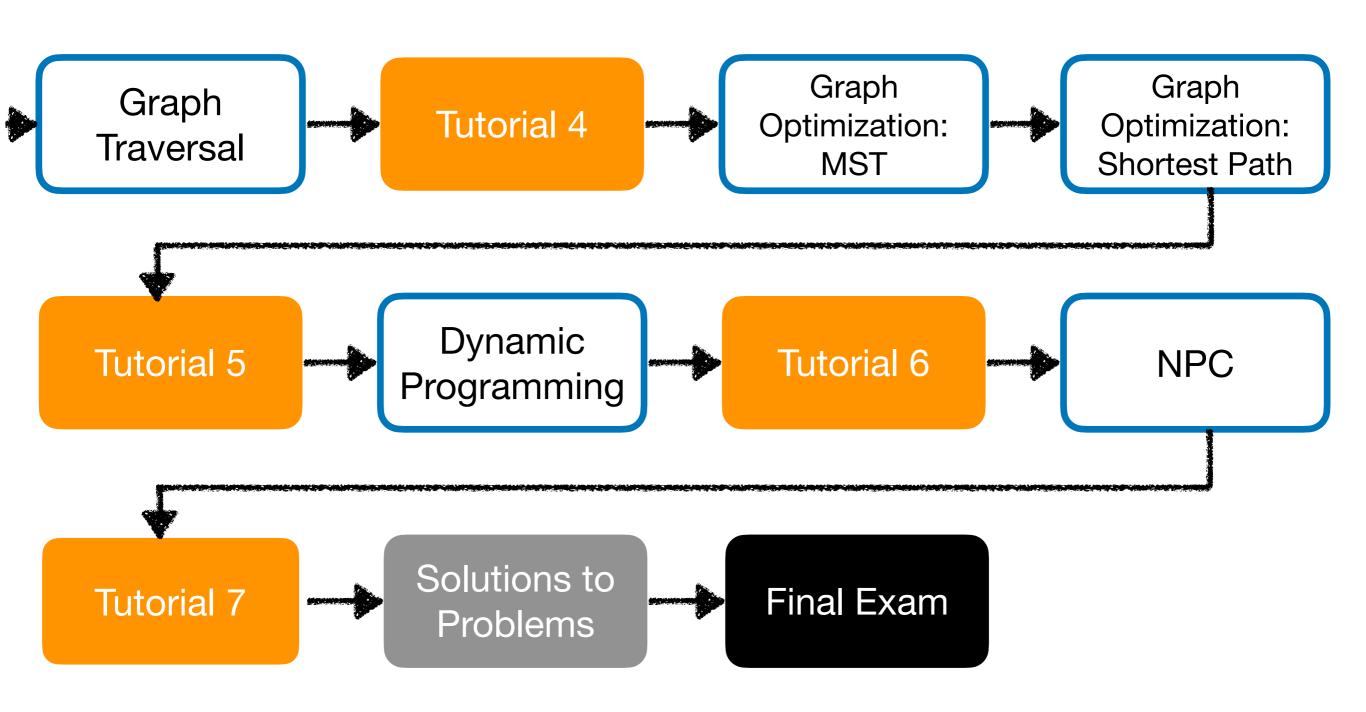
- Syllabus
- Textbook
- Website

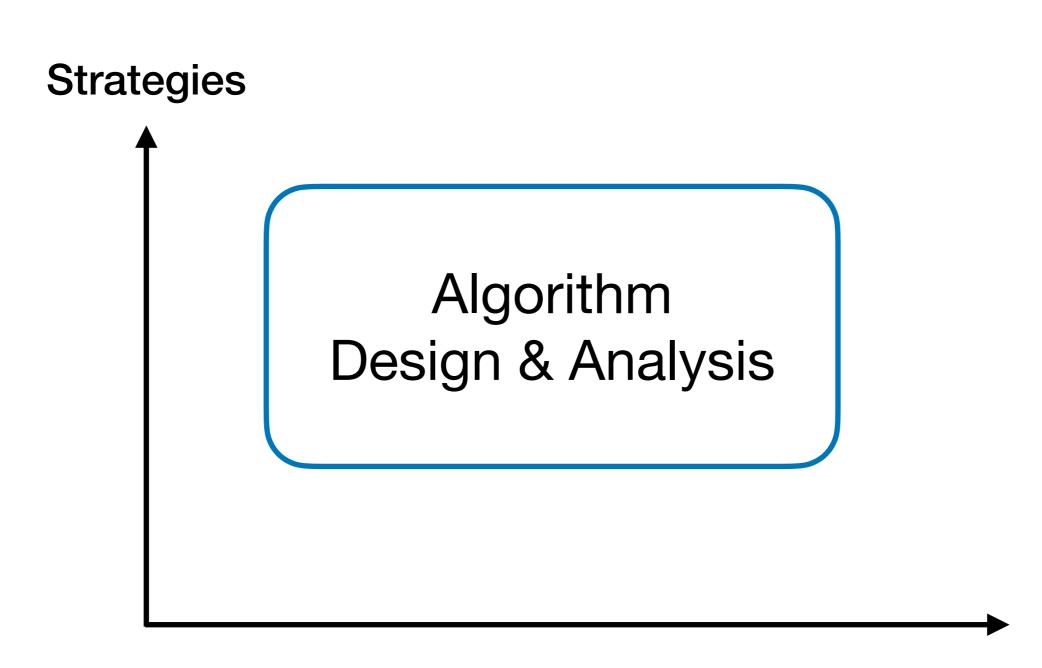
Model of Computation

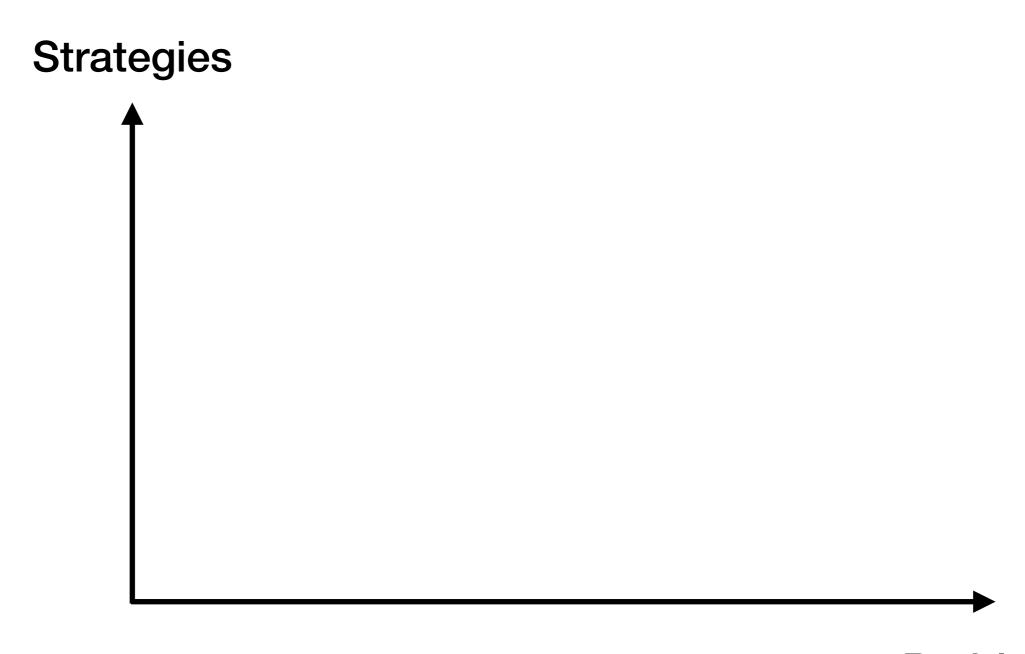
Algorithm
Design &
Analysis
Techniques

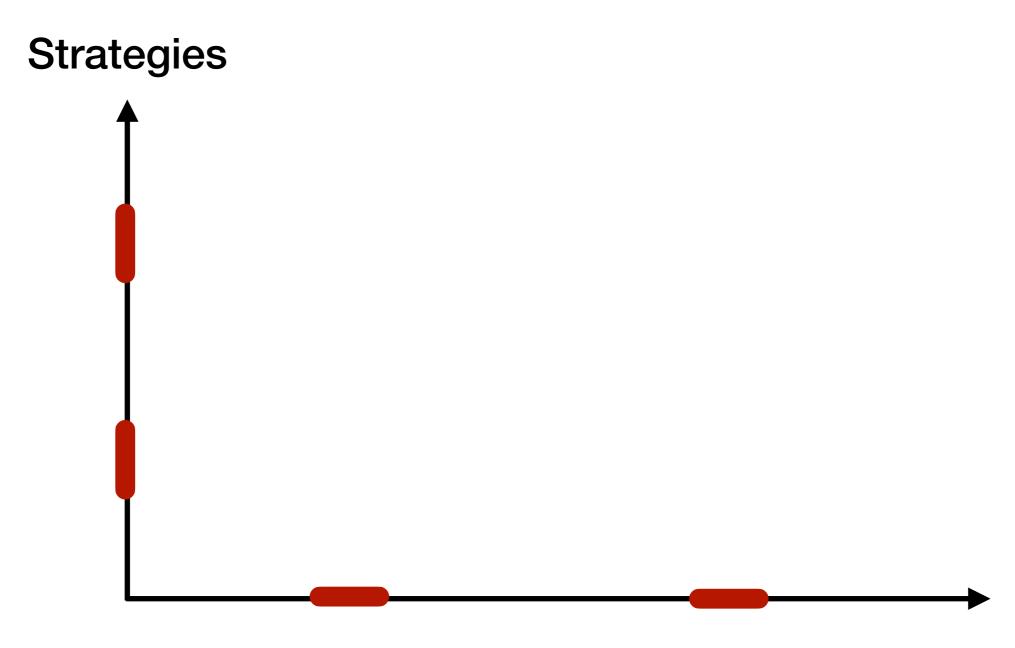
Computation Complexity

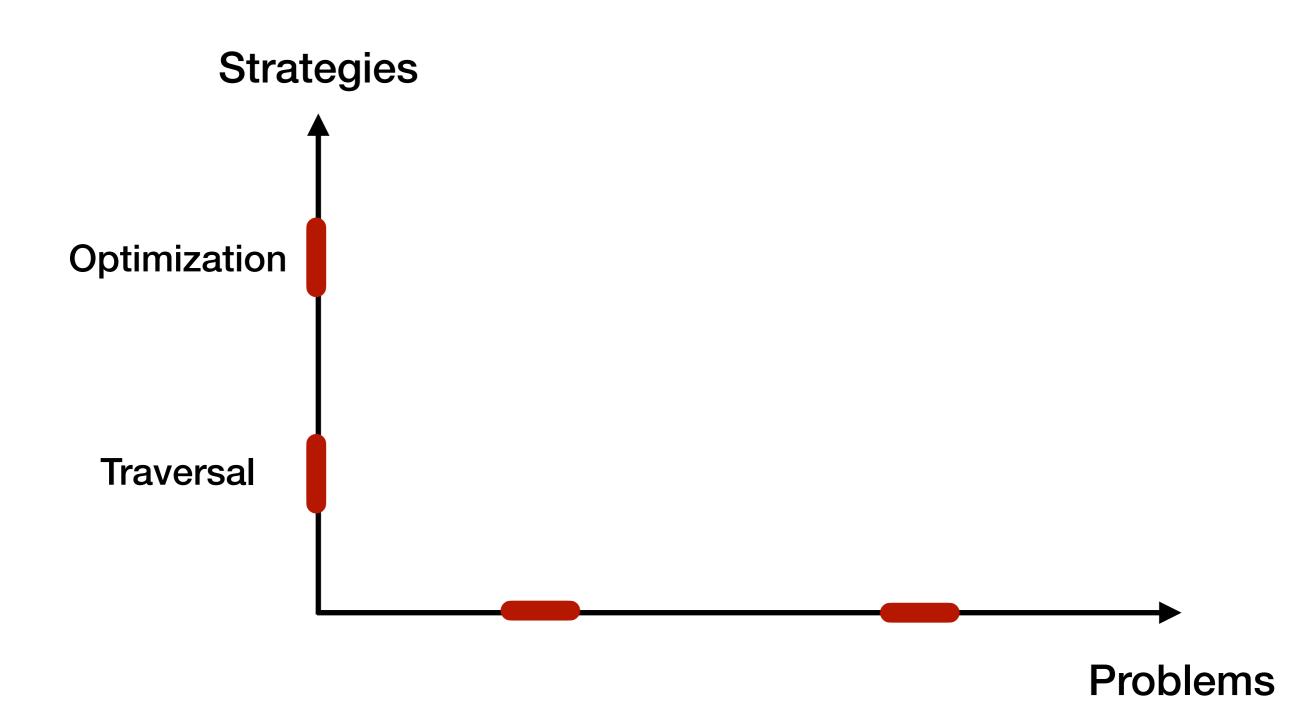


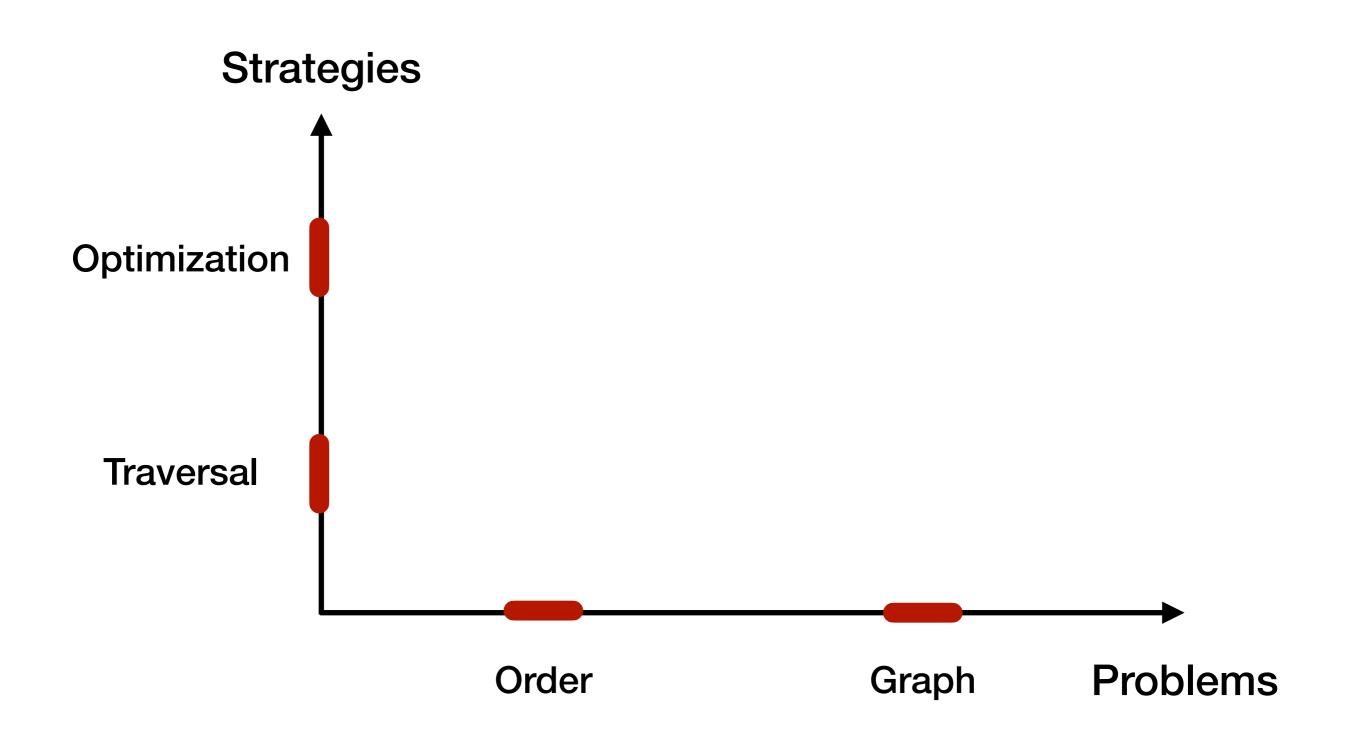


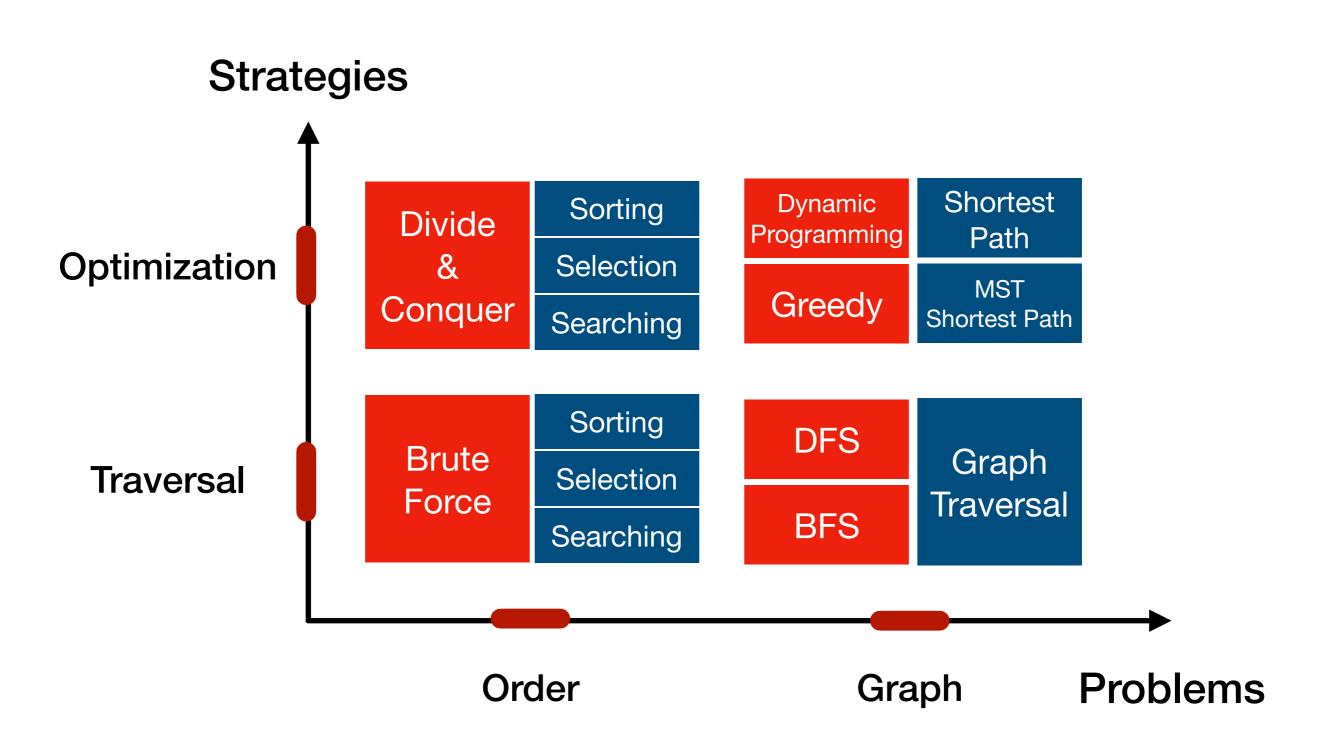






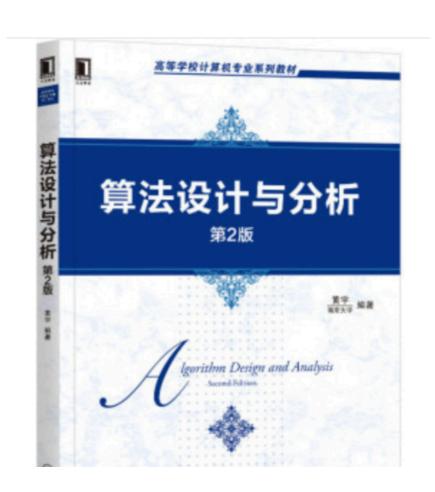






Textbooks

- Course outline: LADA
 - Lectures on Algorithm Design & Analysis (slides)
- Course contents
 - Algorithm Design and Analysis



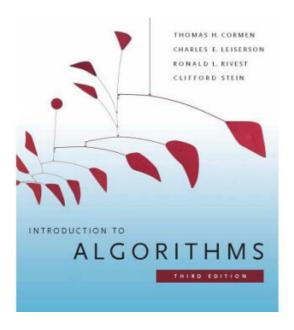
More info about this book: https://zhuanlan.zhihu.com/p/24150569

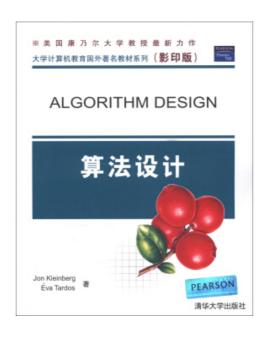
Textbooks

- Further reading
 - Introduction to Algorithms (CLRS)

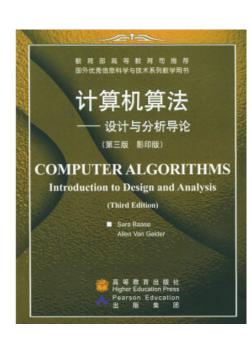
https://walkccc.me/CLRS/

- Algorithms
- Algorithm Design
- Computer Algorithms*









我应该如何学习南京大学徐经纬老师的算法设计与分析课程?

I'm considering the user's request to learn Professor Xu Jingwei's algorithm design and analysis course from Nanjing University, and I'm mapping out the best approaches to provide useful recommendations.

Considering how to respond to the user's request, I am focusing on providing practical and effective advice.

要学习南京大学徐经纬老师的《算法设计与分析》课程, 您可以按照以下步骤进行:

- 1. **获取课程资料**:南京大学计算机科学与技术系已将《算法设计与分析》课程的讲义和幻灯片公开,您可以在系网站上找到相关资料。 CS.NJU.EDU.CN
- 2. **观看视频讲解**:在B站上,您可以找到徐经纬老师的课程视频。例如,2024年南京大学计算机科学与技术系的《算法设计与分析》教程视频已上传至B站。 BILIBILI.COM
- 3. **参考其他优质课程**:在B站上,有许多高质量的计算机科学课程可供参考。例如,东北大学 2021年春季学期的《算法设计与分析》课程视频也已上传至B站。 V2EX.COM
- 4. **参与讨论和实践**:加入相关的学习社区或论坛,与其他学习者交流,讨论课程内容,解决疑问。此外,积极参与编程实践,巩固所学知识。

通过以上步骤,您可以系统地学习徐经纬老师的《算注设计与分析》课程,深入理解算法的设计与分析方法。

Course Websites

- 算法QQ群
 - 1025338112
- Online Judge平台
 - Coming soon



Criteria of Assessment

- 作业 20%
 - 纸质作业
 - OJ平台
- 期中考试 20%
- 期末考试 60%

Algorithm Design & Analysis

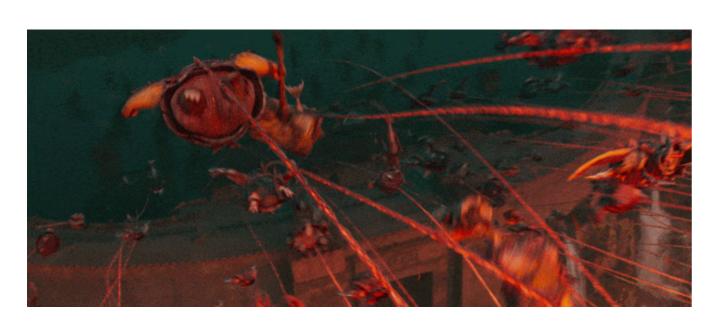
- Algorithm the spirit of computing
 - Model of computation
- Algorithm by example
 - Greatest common divisor
 - Sequential search

Algorithm Design & Analysis

- Algorithm design & analysis
 - Design: correctness
 - Analysis: worst-case / average-case cost

Computer and Computing

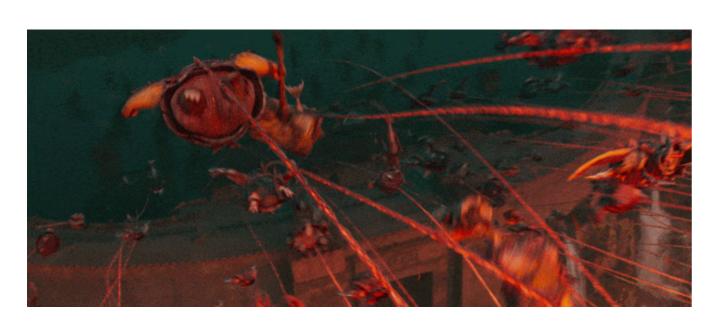
- Problem 1: why the computer seems to be able to do anything?
 - Scientific computing, document processing, computer vision, movies, computer games, ...





Computer and Computing

- Problem 1: why the computer seems to be able to do anything?
 - Scientific computing, document processing, computer vision, movies, computer games, ...





Reasoned for a couple of seconds >

I'm curious about achieving a GPA of 5.0 at Nanjing University's School of Computer Science. This involves deep analysis, strategic planning, and consistent effort, as consistent top scores indicate thorough understanding and mastery of the coursework.

To achieve a GPA of 5.0 at Nanjing University's School of Computer Science, I'll delve into effective study habits, staying organized, maintaining motivation, and seeking guidance from instructors and fellow students.

要在南京大学计算机学院获得 GPA 5.0 的高分,除了需要扎实的学术基础外,还需要良好的学习策略和时间管理能力。以下是一些建议,帮助你在学术生涯中取得优异成绩:

1. 基础知识扎实:

确保你对课程的基础概念和知识点有深入的理解。计算机科学的许多课程都建立在前期的知识基础上,打好基础是关键。

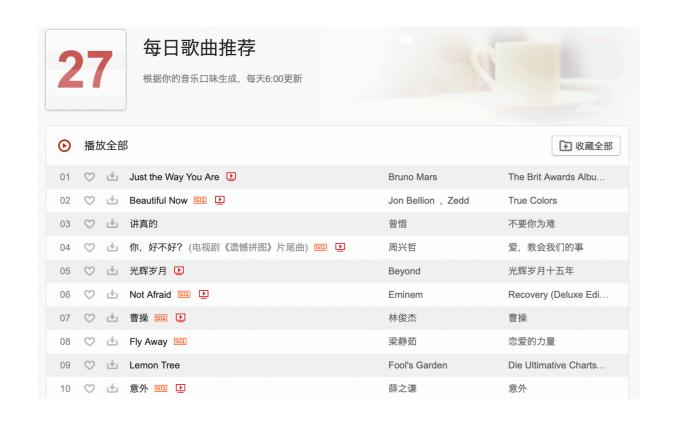
2. 课前预习:

• 课前预习能够帮助你提前了解课程内容和难点,课堂上更容易跟上教授的讲解,提高学习效率。

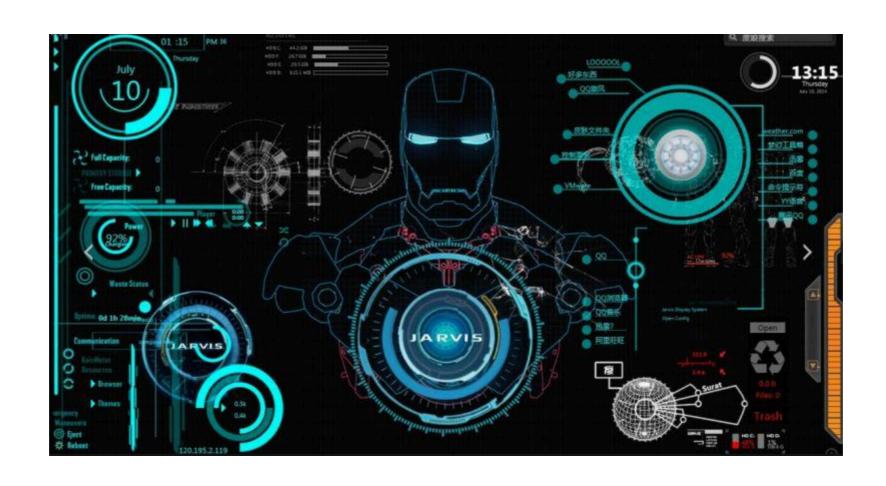
3. 主动参与课堂讨论:

Computer and Computing

- Problem 2: what can / cannot be efficiently done by a computer?
 - Manage millions of songs v.s. music composition



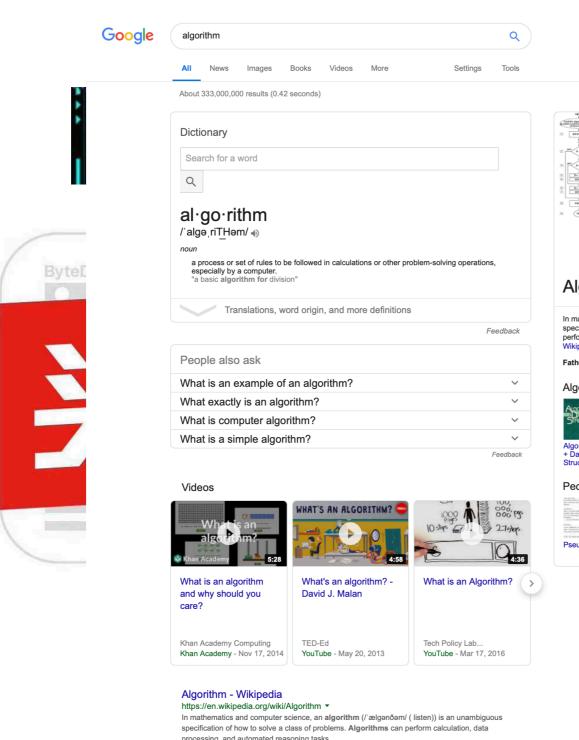


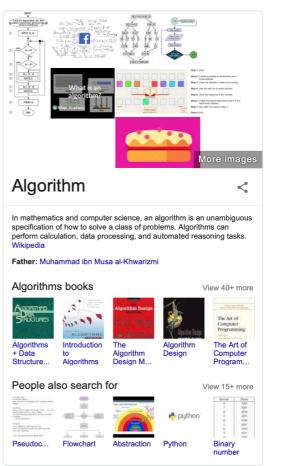










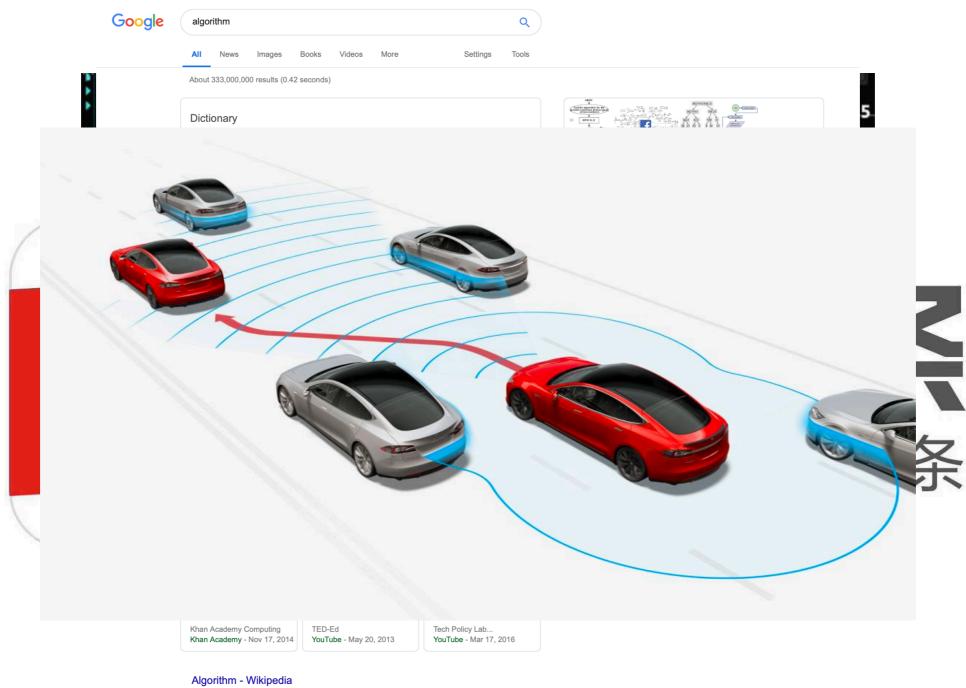


Feedback



processing, and automated reasoning tasks.

Graph algorithms · Algorithm (disambiguation) · Euclidean algorithm · Search



https://en.wikipedia.org/wiki/Algorithm ▼

In mathematics and computer science, an algorithm (/ˈælgenðem/ (listen)) is an unambiguous specification of how to solve a class of problems. Algorithms can perform calculation, data processing, and automated reasoning tasks.

 $\textbf{Graph algorithm} \cdot \textbf{Algorithm (disambiguation)} \cdot \textbf{Euclidean algorithm} \cdot \textbf{Search}$

🎉 DeepSeek-R1 已发布并开源,性能对标 OpenAl o1 正式版,在网页端、APP 和 API 全面上线,点击查看详情。

deepsek

探索未至之境

开始对话

免费与 DeepSeek-V3 对话 使用全新旗舰模型

获取手机 App

DeepSeek 官方推出的免费 AI 助手 搜索写作阅读解题翻译工具

Khan Academy Computing Khan Academy - Nov 17, 2014 YouTube - May 20, 2013

YouTube - Mar 17, 2016

Algorithm - Wikipedia

https://en.wikipedia.org/wiki/Algorithm •

In mathematics and computer science, an algorithm (/ˈælgənðəm/ (listen)) is an unambiguous specification of how to solve a class of problems. Algorithms can perform calculation, data processing, and automated reasoning tasks.

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Algorithm





Algorithm





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探索未至之境

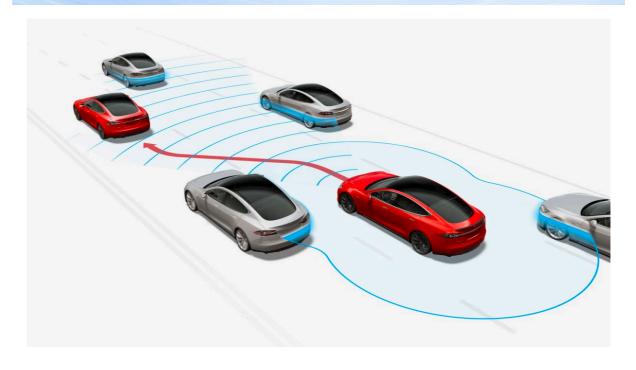
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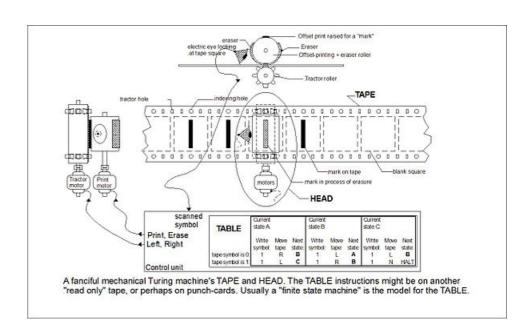


Computer and Computing

Computing

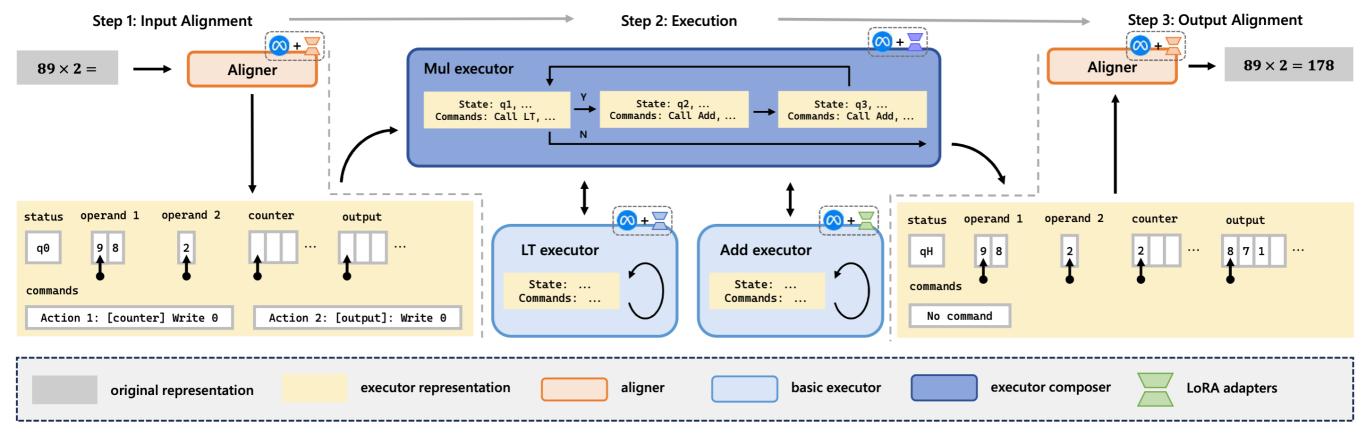
- Encoding everything into '0's and '1's
- Operations over '1's and '0's
- Decoding the '1's and '0's
- Turing machine
 - An abstract/logical computer





Executing Arithmetic: Fine-Tuning Large Language Models as Turing Machines

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{caochun,jingweix}@nju.edu.cn



Algorithm

- Algorithm is the spirit of computing
 - To solve a specific problem (so-called an algorithmic problem)
 - Combination of basic operations
 - in a precise and elegant way



Algorithm

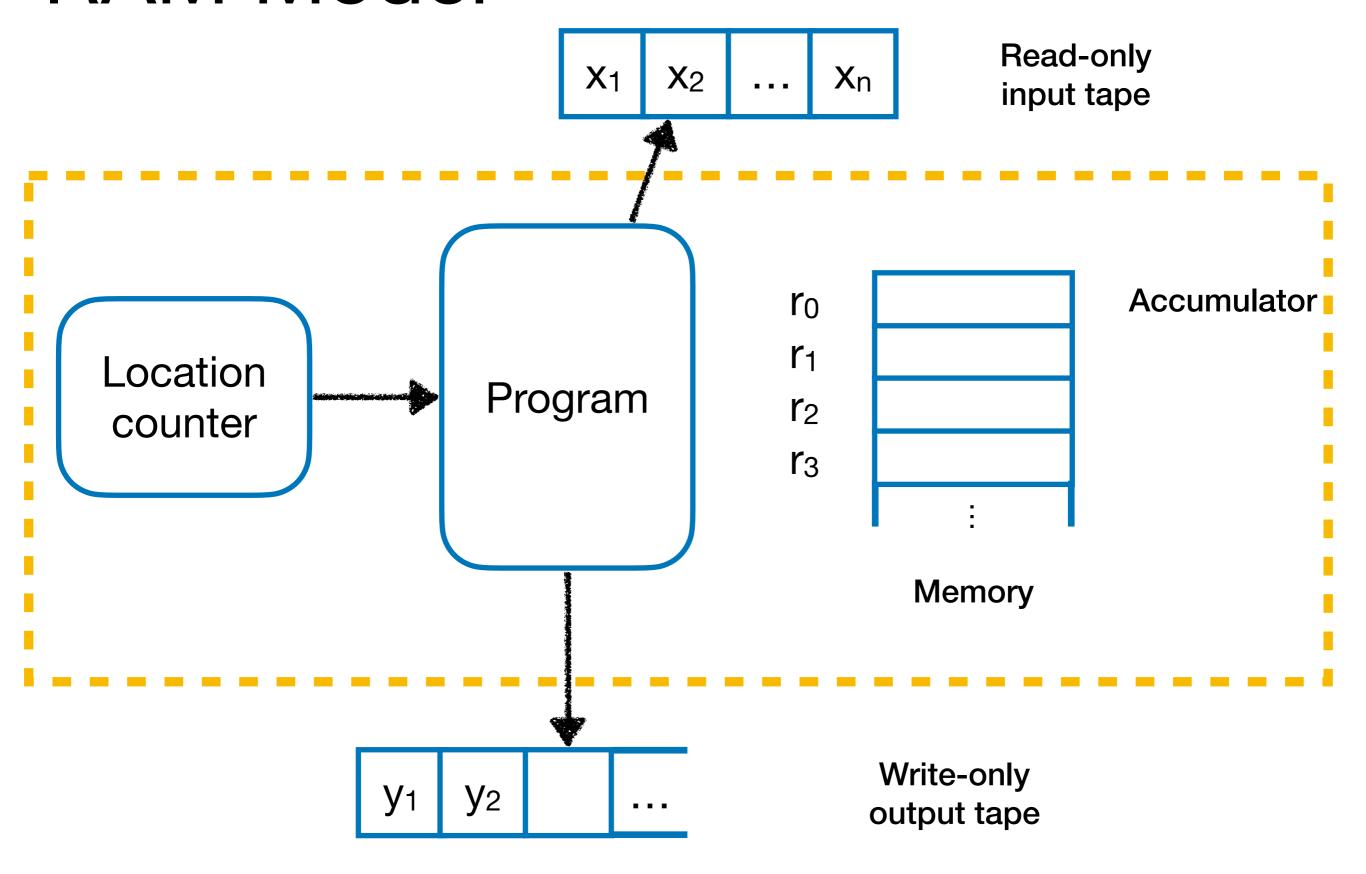
- Essential issues
 - Model of computation
 - Algorithm design
 - Algorithm analysis

Model of Computation

Problems

- Why the algorithms we learn can run almost everywhere?
- Why the algorithms we learn can be implemented in any language?
- Machine- and language- independent algorithms, running on an abstract machine
 - Turing machine: over-qualify
 - RAM model: simple but powerful

RAM Model



The RAM Model of Computation

- Each simple operation takes one time step
 - E.g., key comparison, +/-, memory access, ...
- Non-simple operations should be decomposed
 - Loop
 - Subroutine

- Memory
 - Memory access is simple operation
 - Unlimited memory

Tradeoff: accuracy v.s. ease of use

To Create an Algorithm

Algorithm design

Composition of simple operations, to solve an algorithmic problem

- Amount of work done / memory used
 - In the worst/average case
- Advanced issues
 - Optimality, approximation ratio, ...

Algorithm by Example

- Algorithmic Problem 1
 - Find the greatest common divisor of two nonnegative integers m and n
- Algorithmic Problem 2
 - Is a specific key K stored in array E[1..n]?

Probably the Oldest Algorithm

Euclid Algorithm

 Find the greatest common divisor of two non-negative integers m and n

Specification

Input: non-negative integer m, n

Output: gcd(m, n)

Euclid algorithm

[E1] n divides m, the remainder -> r

[E2] if r=0 then return n

[E3] n -> m; r -> n; goto E1

Euclid algorithm - recursive version

Euclid(m, n)

[E1] if r=0 then return m

[E2] else return Euclid(n, m mod n)

Sequential Search

Problem

 Search an array for a specific key

Specification

```
Input: K, E[1..n]
Output: Location of K (1,2,...,n; -1: K is not in E[])
```

Sequential search

```
int seqSearch(int[] E, int n, int K)
  int ans, index;
  ans = -1;
  for(index = 1; index <= n; index ++)
     if(K==E[index])
     ans = index;
     break;
  return ans;</pre>
```

Algorithm Design

- Criteria
 - Defining correctness
- Main challenge
 - For proving correctness
- Our strategy
 - Mathematical induction
 - ...

Specification

Input: non-negative integer m, n

Output: gcd(m, n)

Main challenge

The output is always correct, for any legal input.

Infinite possible inputs

Mathematical induction

Weak principle Strong principle

For Your Reference

Mathematical induction

The Weak Principle of Mathematical Induction

If the statement p(b) is true and the statement p(n-1) => p(n) is true for all n>b, then p(n) is true for all integers n>=b.

The Strong Principle of Mathematical Induction

If the statement p(b) is true, and the statement p(b) and p(b+1) and ... and p(n-1) => p(n) is true, for all n>b, then p(n) is true for all integers n>=b.

Correctness of the Euclid Algorithm

- Induction on n
 - Base case
 - n = 0: for any m, Euclid(m, 0) = m;
 - n = 1: for any m, Euclid(m, 1) = 1;
 - n = 2: ...
 - Assumption
 - For any n≤N₀, Euclid(m, n) is correct;
 - Induction
 - Euclid(m, N₀+1) = Euclid(N₀+1, m mod (N₀+1));

Notes on Induction

"Notes on Structured Programming", E.W. Dijkstra

I have mentioned mathematical induction explicitly, because it is the only pattern of reasoning that I am aware of, that eventually enables us to cope with loops and recursive procedures.

- Criteria
 - Performance metrics
- Worst case
 - Best case?
- Average case
 - Average cost?
- Advanced topics
 - Lower bound, optimality, ...

- How to measure
 - Not too general
 - Giving essential indication in comparison of algorithms
 - Not too precise
 - Machine independent
 - Language independent
 - Programming paradigm independent
 - Implementation independent

Criteria

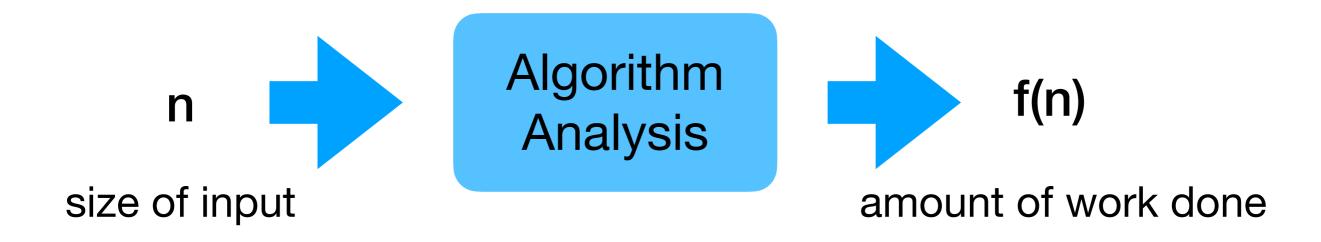
Critical operation

How many critical operation are conducted

For example

Algorithmic problem	Critical operation
Sorting, selection, searching String matching	Comparison (of keys)
Graph traversal	Processing a node/edge
Matrix multiplication	Multiplication

- Amount of work done
 - usually depends on size of the input
 - usually does not depend on size of the input only



Worst-case Complexity

- W(n)
 - Upper bound of cost
 - For any possible input
 - $W(n) = \max_{I \in D_n} f(I)$

Average-case Complexity

- A(n)
 - Weighted average

$$A(n) = \sum_{I \in D(n)} Pr(I)f(I)$$

- A special case
 - Average cost
 - Total cost of all inputs, averaged over the input size

•
$$A(n) = \frac{1}{|D(n)|} \sum_{I \in D(n)} f(I)$$

Average-case Cost of SeqSearch

- Case 1: K is in E[]
 - Assumptions:
 - Assuming that K is in E[]
 - Assuming no same entries in E[]
 - Each possible input appears with equality (thus, K in the ith location with probability 1/n)

$$A_{succ}(n) = \sum_{i=0}^{n-1} Pr(I_i|succ)t(I_i) = \sum_{I=0}^{n-1} \frac{1}{n}(i+1) = \frac{n+1}{2}$$

Average-case Cost of SeqSearch

- Case 2: K may (or may not) be in E[]
 - Assume that K is in E[] with probability q

$$A(n) = Pr(succ)A_{succ}(n) + Pr(fail)A_{fail}(n)$$
$$= q\frac{n+1}{2} + (1-q)n$$

How to make reasonable assumptions?

- Advanced topics
 - Lower bound (Selection)
 - Optimality (Greedy, DP)
 - Computation complexity
 - Approximate / online / randomized algorithms

Thank you! Q & A